



# Optimizing Floating Point Calculations, I

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# Overview

- Introduction
- Computer Basics
- Intel<sup>®</sup> Pentium<sup>®</sup> 4 Hardware Architecture
- Memory Hierarchy
- Cache Basics
- Measuring Program Performance

# Introduction

- Focus on writing computer programs in C, FORTRAN
- UNIX (Linux) platform
- Intel<sup>®</sup> Pentium<sup>®</sup> Hardware Platform
- While above seem very specific, concepts and techniques are applicable to most programming environments on most hardware platforms.

# Computer Basics

- The Five Classic Components of a Computer:
  - Input
  - Output
  - Memory
  - Datapath
  - Control
- Datapath and Control are the domain of the *Central Processing Unit* (CPU) or Processor.
- The CPU runs at a specified *clock rate* which relates to how fast the hardware can perform basic functions. Pentium 4 speeds 3 - 4 GHz.
- Memory access speed typically ranges 400 - 800 MHz.

# Computer Basics, contd.

- Associated with each CPU is an *instruction set* which define what actions the CPU can take.
- A CPU can be classified as either a *Reduced Instruction Set Computer (RISC)* or a *Complex Instruction Set Computer (CISC)*.
- A RISC computer can recognize and execute only a small number of instructions, usually  $< 128$ .
- A CISC computer has a much larger set of instructions it can recognize and execute.
- Modern Intel<sup>®</sup> CPU's are CISC in principle but utilize RISC concepts in implementation.

# Computer Basics, contd.

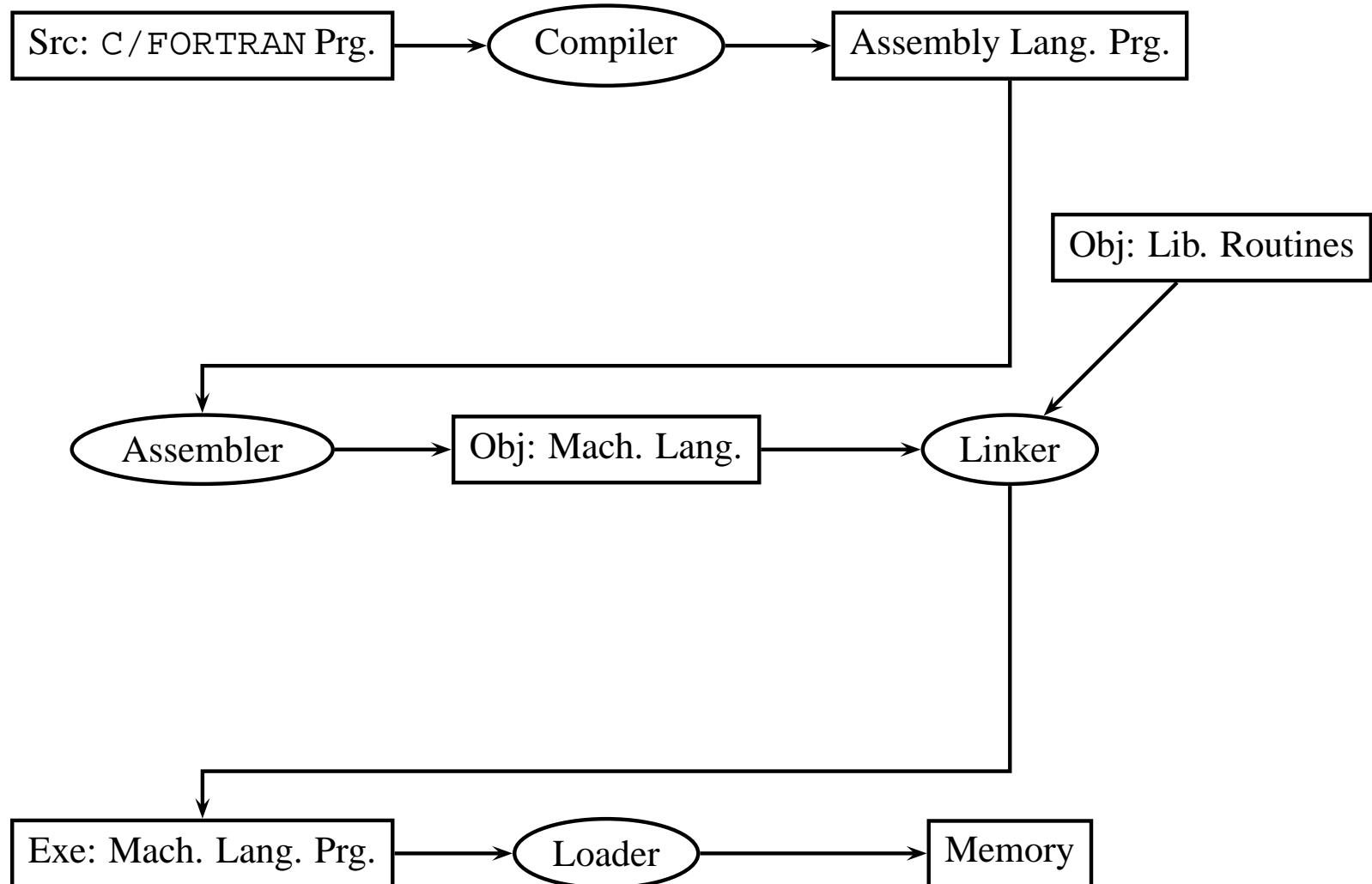
- Today's modern CPU's consist of a number of different components:
  - Fetch/Decode Unit - Fetches instructions to be executed and Decodes into instructions which can be processed.
  - Arithmetic Logical Unit (ALU) - consists of specialized processor functions for dealing with operations involving integer data, floating point data, etc.
  - Write Unit - Handles output to memory.
- *Pipelining* is a technique in which multiple instructions are overlapped in execution.

# Computer Basics, contd.

- The bottom line is that a CPU acts on a stream of ones and zeros which it interprets to perform specific operations.
- A *executable program* is simply this binary stream.
- In order to allow one to write a program of any complexity, we have developed a method to translate commands from a *high level language* to *assembly language* to *machine language* which is processed by the CPU.
- This process of translation is called the *Compilation Process*.

# The Compilation Process

High level program must be translated into machine language and loaded into memory to run.





# Binary Basics

- All CPU's have located on the chip a small number of storage locations called *registers*.
- These registers are described by how many bits they contain.
- A 32 bit processor implies that the CPU acts mostly on registers containing 32 bits.
- A byte is 8 bits.
- A 32 bit address space implies that there are  $2^{32}$  uniquely addressable bytes.

# Intel<sup>®</sup> CPU History

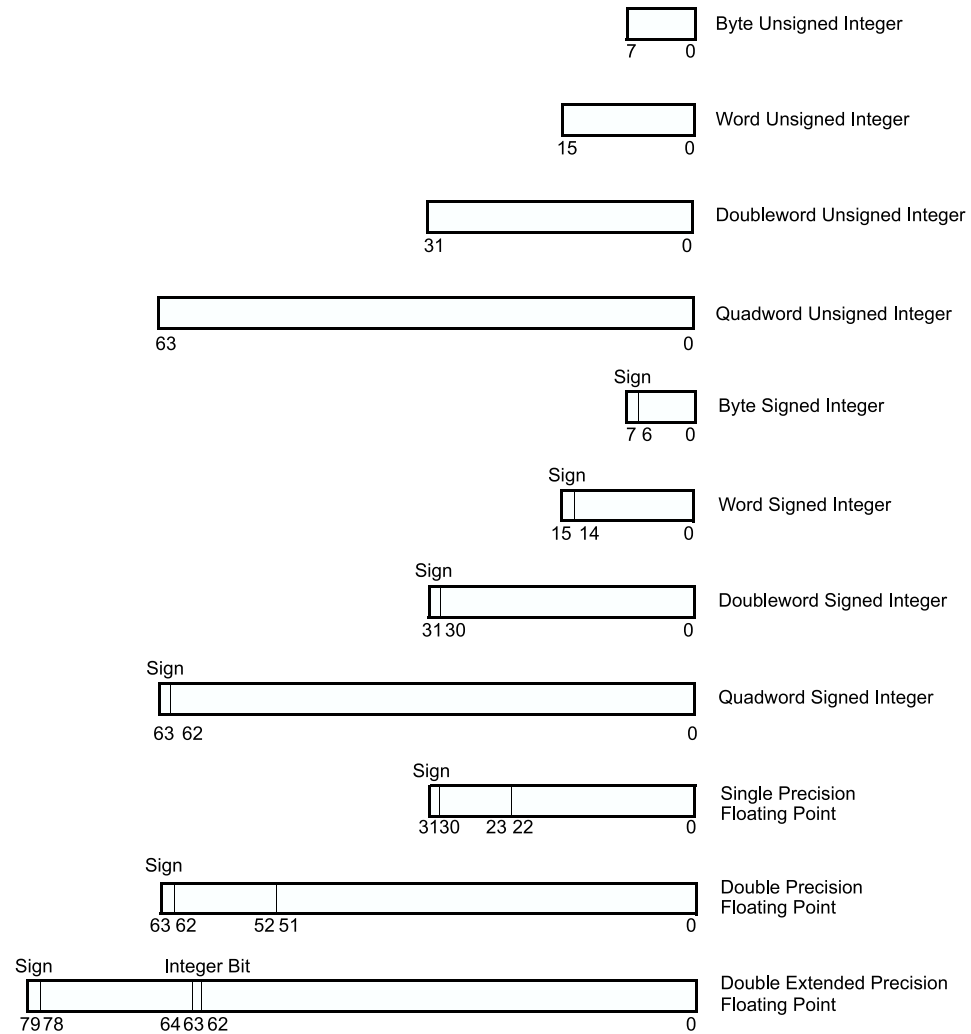
- 1978** 8086 arch., extension of 8 bit microprocessor 8080. 16 bit arch., internal registers 16 bits wide.
- 1980** 8087 Floating Point Coprocessor extends 8086 instruction set with ~ 60 floating point instructions. Stack based instead of register based.
- 1982** 80286 increases address space to 24 bits, elaborate memory mapping model.
- 1985** 80386 extends address space to 32 bits, more flexible register usage.
- 1989** 80486.
- 1992** Pentium.
- 1995** Pentium Pro.
- 1997** MMX extensions. 57 new instructions using floating point stack.
- 1999** Pentium III. SSE (Streaming Single Instruction Multiple Data [SIMD]) extensions. 70 new instructions. Added 8 separate registers, doubled width to 128 bits, SP packed data type implies that 4 32 bit FP Ops can be done in parallel. Also includes cache prefetch and streaming store instructions which bypass cache and write directly to memory.

# Intel<sup>®</sup> CPU History, contd.

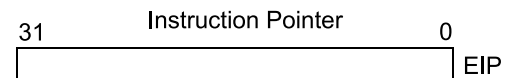
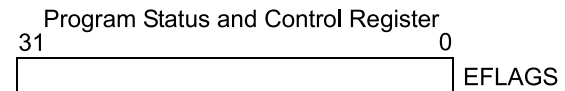
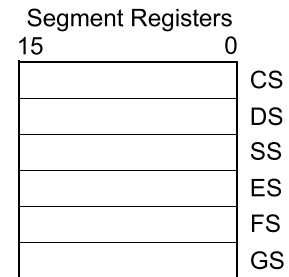
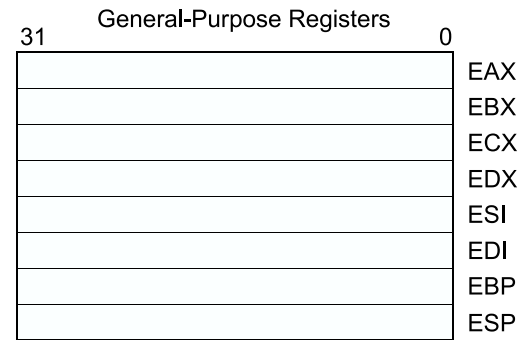
- 2001** Pentium 4 with SSE2 extensions. 144 new instructions. DP packed data type implies that 2 64 bit FP Ops can be done in parallel. Compilers can choose to use 8 SSE registers as FP registers.
- 2003** AMD64 extends address space from 32 to 64 bits. Increases number of registers to 16 and 128 bit registers to 16.
- 2004** Intel embraces AMD64 extensions producing EM64T. SSE3 adds 13 instructions to support complex arithmetic.

*IA-32* is commonly used to describe Intel Pentium 4 CPU's. These processors have 8 general purpose registers (GPR) each 32 bits wide. Note that this is the same since 80386 introduced in 1985.

















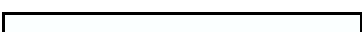
# Numeric Datatypes



# General Purpose Registers



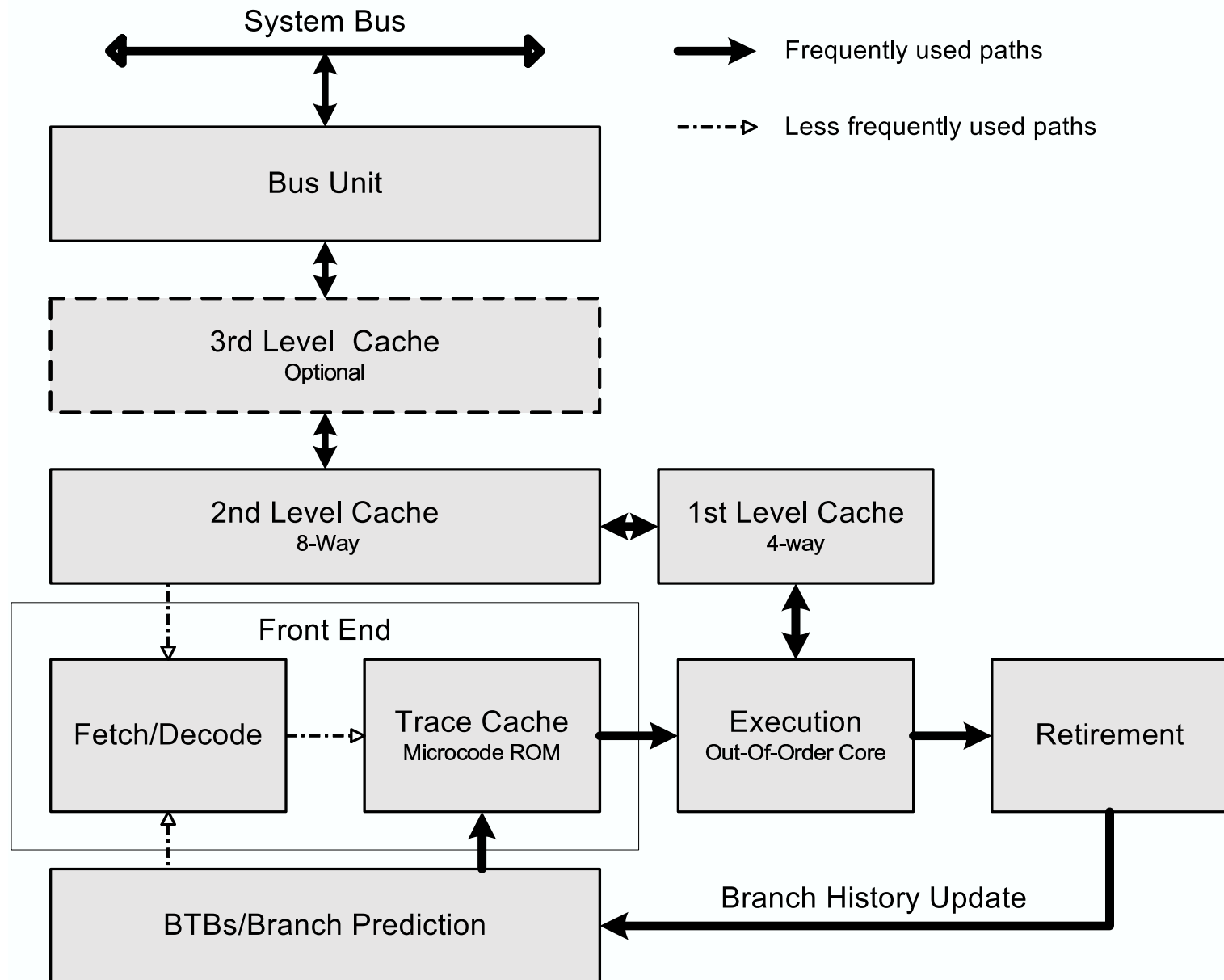
# SIMD Registers

SIMD Extension	Register Layout	Data Type
MMX Technology	MMX Registers 	8 Packed Byte Integers
		4 Packed Word Integers
		2 Packed Doubleword Integers
		Quadword
SSE	MMX Registers 	8 Packed Byte Integers
		4 Packed Word Integers
		2 Packed Doubleword Integers
		Quadword
	XMM Registers 	4 Packed Single-Precision Floating-Point Values
SSE2/SSE3	MMX Registers 	2 Packed Doubleword Integers
		Quadword
	XMM Registers 	2 Packed Double-Precision Floating-Point Values
		16 Packed Byte Integers
		8 Packed Word Integers
		4 Packed Doubleword Integers
		2 Quadword Integers
	Double Quadword	

# IA-32 Register Summary

- Linear address space 32 bit (4 GB), physical address space 36 bit (64 GB)
- 8 GPR, 6 segment registers (addressing), EFLAGS register (control), EIP register (instruction pointer)
- 8 x87 FPU data registers, misc FPU control registers
- 8 MMX registers for SIMD operations on 64-bit packed byte, word, doubleword integers
- 8 XMM data registers for SIMD operations on 128-bit packed SP and DP FP data and on 128-bit packed byte, word, doubleword, and quadword integers.

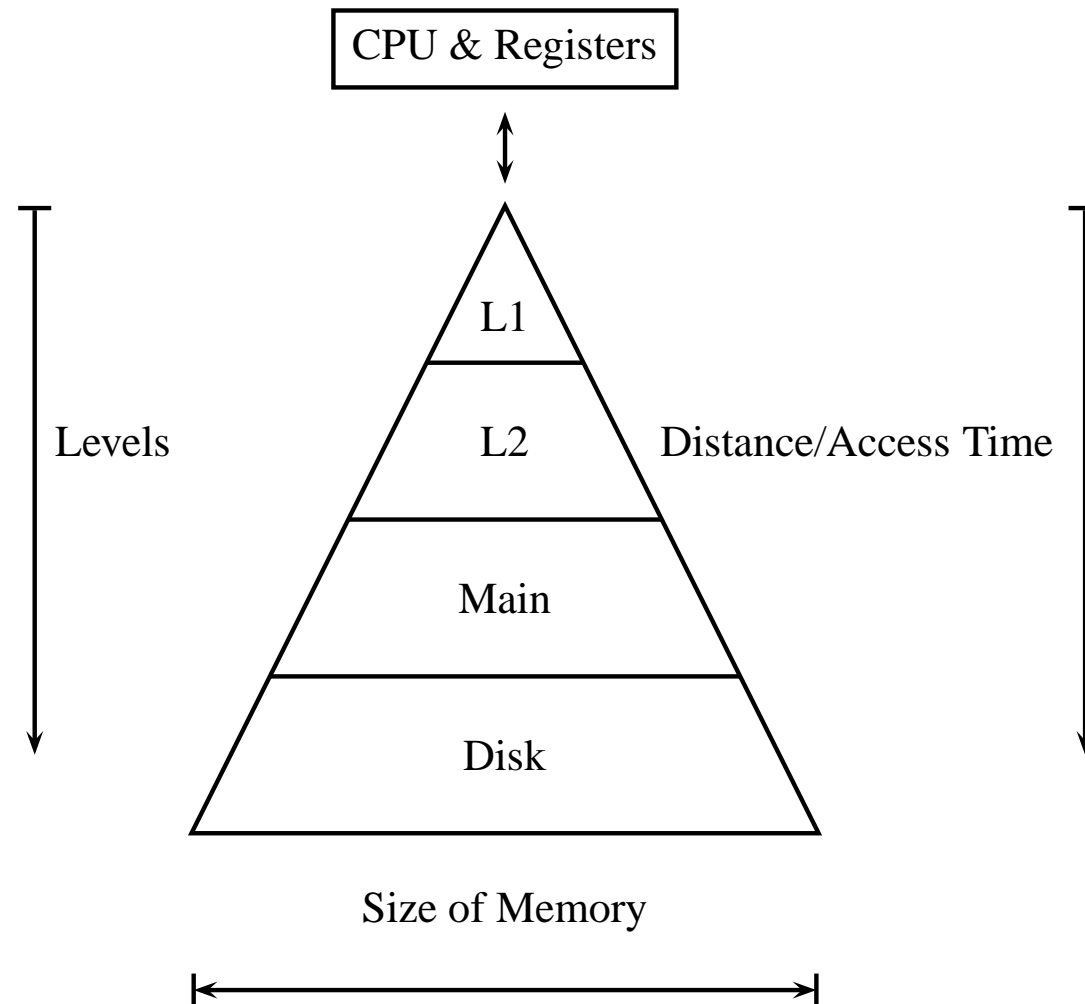
# Intel® Netburst® Architecture





# Cache Basics

Memory on a computer is organized in a hierarchical manner:



# Cache Basics, contd.

- Memory Characteristics

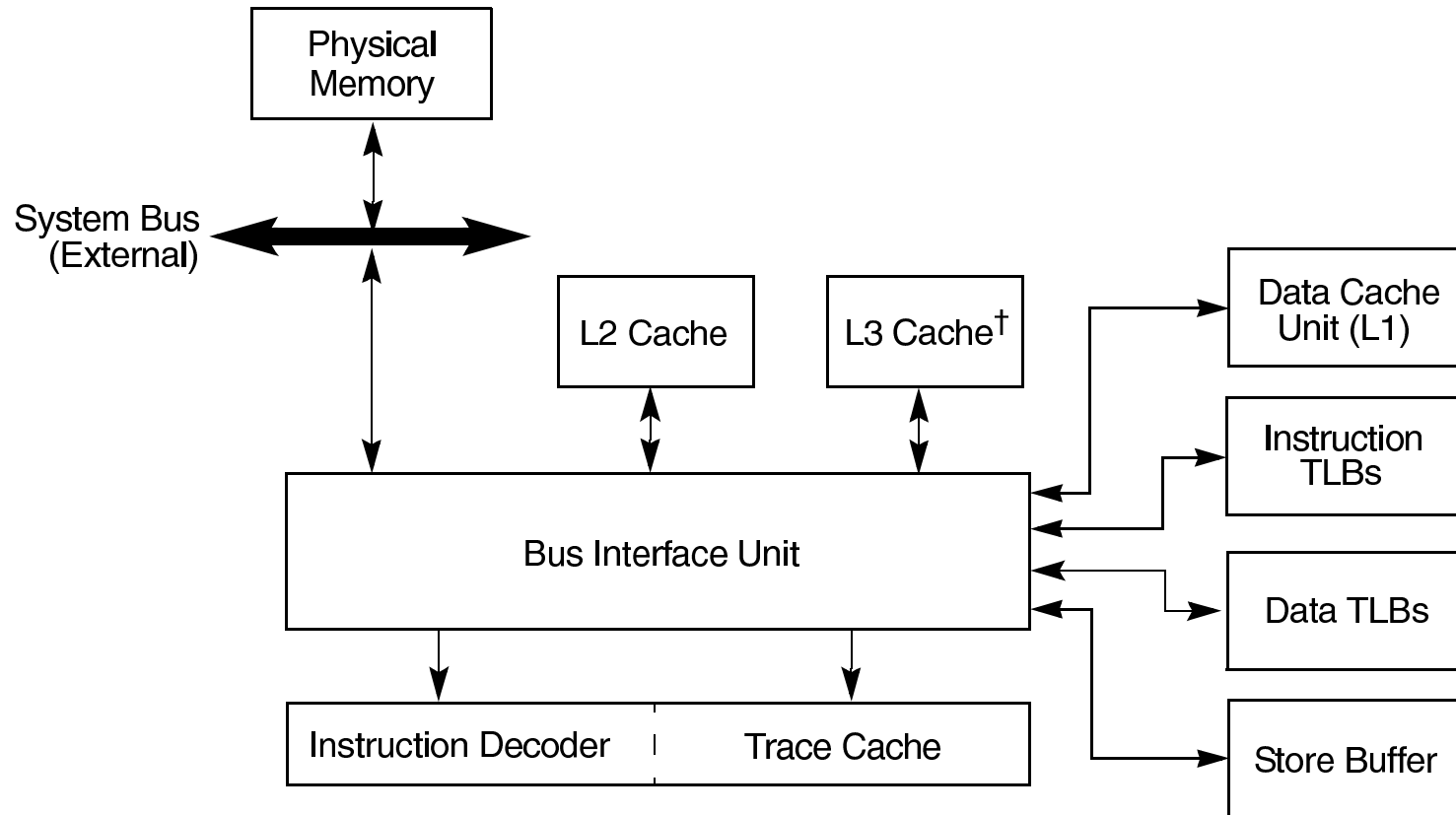
2004 \$/GB	Type	Access Time (ns)
\$4K - \$10K	SRAM - Static (On Chip Caches)	0.5 - 5
\$100 - \$200	DRAM - Dynamic (Main Memory)	50 - 70
\$0.5 - \$2	Disk	5e6 - 20e6

- An additional area of memory is called the *Translation Lookaside Buffer (TLB)* which contains virtual – physical address translations.
- TLB's keep track of large *pages* of memory that are in use, either 4KB, or 2MB/4MB starting addresses.
- When data is needed, best to retrieve from lowest level cache if available.

# Cache Basics, contd.

- *Temporal locality* - Principle that if a data location is referenced it will be referenced again soon.
- *Spatial locality* - Principle that if a data location is referenced, data locations with nearby addresses tend to be referenced soon.
- A cache *line* is usually either 64 or 128 bytes of contiguous storage.
- When CPU requires a data item not currently in cache (L1,L2) it will read an entire line of data into the appropriate cache (L1, L2, or all).
- Levels in cache hierarchy are not inclusive. The fact that a line is in Level  $i$  does NOT imply that it is also in Level  $i + 1$ .

# Cache Structure



† Intel Xeon processors only

# Cache Characteristics

- Cache Parameters

Characteristic	Intel Pentium 4	AMD Opteron
L1 Cache Size	8 KB data, 12K instr trace	64 KB data, 64 KB instr
L1 Assoc.	4-way set assoc.	2-way set assoc.
L1 Line Size	64 bytes	64 bytes
L1 Policy	LRU, Write-through	LRU, Write-back
L2 Cache Size	512 KB (data and instr)	1 MB (data and instr)
L2 Assoc.	8-way set assoc.	16-way set assoc.
L2 Line Size	128 bytes	64 bytes
L2 Policy	LRU, Write-back	LRU, Write-back

- Associativity refers to how many locations in cache a line can be mapped to.
- Policy refers to Least Recently Used (LRU), i.e., what gets replaced.
- Write-through and Write-back refer to ensuring data integrity when writing to cache.

# Program Performance

- The only complete and reliable measure of computer performance is time.

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$$\text{Time} = \frac{\text{sec}}{\text{Prog}} = \frac{\text{Instr}}{\text{Prog}} \times \frac{\text{Clk cycles}}{\text{Instr}} \times \frac{\text{sec}}{\text{Clk cycle}}$$

- Basic components

Component	Unit
CPU execution time	seconds
Instr Count	# Instr retired
Clk cycles/Instr (CPI)	Avg. # Clk cycles/Instr
Clk cycle time	seconds/Clk cycle

# Program Performance, contd.

- Algorithm affects Instr Count, possibly CPI
- Programming Language affects Instr Count, CPI
- Compiler affects Instr Count, CPI
- Instr Set Arch. affects Instr Count, clock rate, CPI
- Significant program performance degradation can occur when cache misses (including TLB misses) occur on a frequent basis.

# Conclusions

- Intel<sup>®</sup> Pentium<sup>®</sup> 4 has much better computational capabilities than earlier Intel processors.
- Pipelined execution on Pentium 4 allows up to 3 IA-32 instr to execute in a single clock cycle.
- Pentium 4 L1 Cache is small.
- Pentium 4 MMX/SSE/SSE2/SSE3 extensions provide enhanced computational capabilities and allow operating on multiple data items at one time (SIMD).
- To achieve optimum floating point calculation performance, need to be able to use these capabilities.



# Next Week, Part II

- Algorithms and techniques to utilize cache better.
- Compiler choices and options to enable utilization of enhanced features.
- High performance libraries designed to run fast, whats available and how to use.
- Performance Monitoring of applications and interpretation of results.
- Post-processing and visualization of data and results.