

# Communcation over interference channels

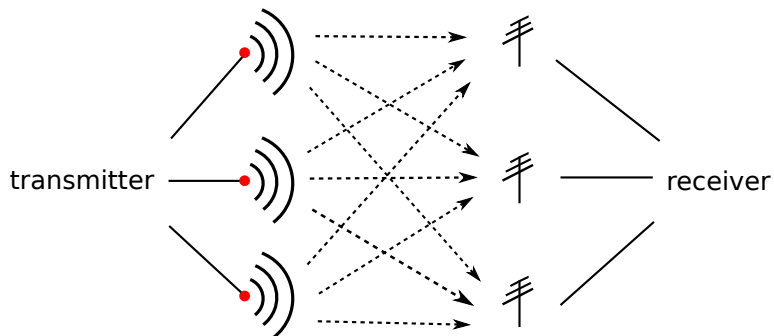
Dustin Cartwright<sup>1</sup>

February 24, 2011

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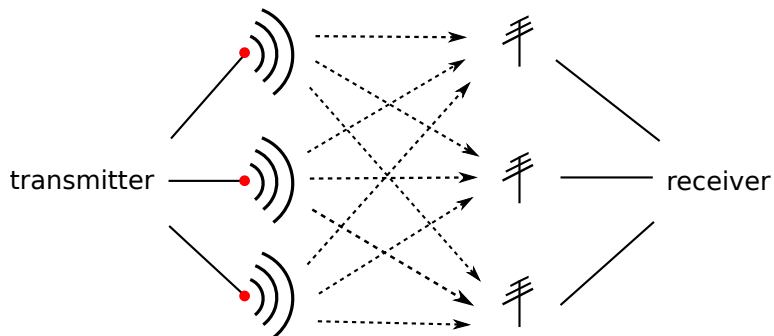
<sup>1</sup>work in progress with Guy Bresler and David Tse

## Multiple-input multiple-output channel



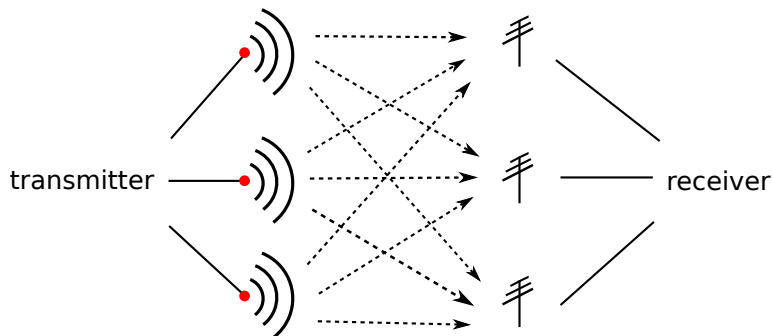
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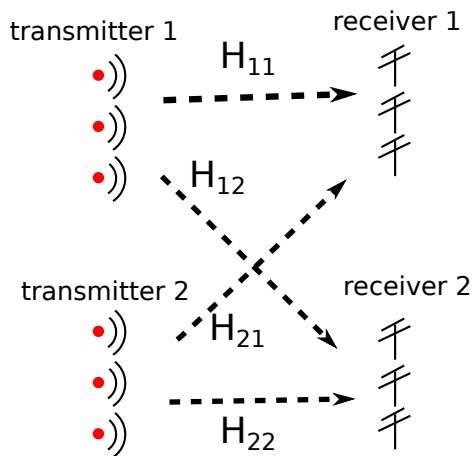
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- ▶ The receiver detects  $Hv \in \mathbb{C}^N$  across its  $N$  antennas, where each entry of the  $H \in \mathbb{C}^{N \times N}$  depends on the signal path.

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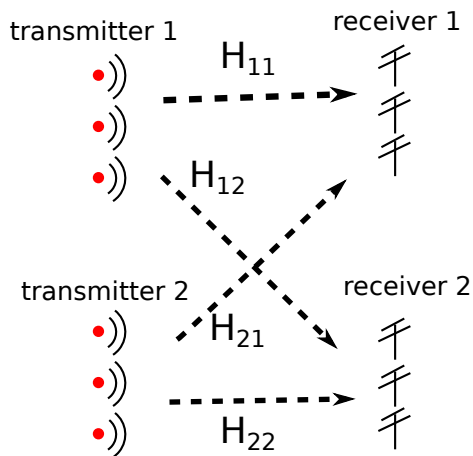
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- ▶ The receiver detects  $Hv \in \mathbb{C}^N$  across its  $N$  antennas, where each entry of the  $H \in \mathbb{C}^{N \times N}$  depends on the signal path.
- ▶ If  $H$  is known and invertible, then the receiver can reconstruct the message  $v$ .

## Multiple users of the same channel



- ▶  $K$  transmitter-receiver pairs using the **same channel**.
- ▶ Determined by  $K^2$  **channel matrices**  $H_{ij}$  of size  $N \times N$ .

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- ▶ Receiver 1 only cares about transmitter 1's message, etc.

## Strategies for interference alignment

- ▶ Each transmitter has a subspace  $V_j \subset \mathbb{C}^N$  to transmit in.
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In order for this to work, we need:

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If each  $H_{ij}$  is **generic**, the second condition is satisfied automatically.

## Questions

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- ▶ What is the **information capacity** of this channel?
- ▶ How to parametrize spaces of solution strategies?

## Incidence correspondence

$$\left(\mathbb{C}^{N \times N}\right)^{K(K-1)} \times \prod_{i=1}^K \text{Gr}(d_i, N) \times \text{Gr}(N - d_i, N)$$

Subvariety of those

$$(H_{12}, \dots, H_{K-1,K}, V_1, \dots, V_K, U_1, \dots, U_K)$$

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### Question

Is the projection onto  $\left(\mathbb{C}^{N \times N}\right)^{K(K-1)}$  surjective?

## Existence of solutions

### Theorem

Assume that  $d = d_1 = \dots = d_K$  and  $K \geq 3$ . Then a generic set of channel matrices has a solution *if and only if*

$$2N \geq d(K + 1).$$

If so, the dimension of the solution variety is

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For non-constant  $d_i$ , we have the **necessary** conditions:

$$d_i + d_j \leq N \quad \text{for all } i, j$$

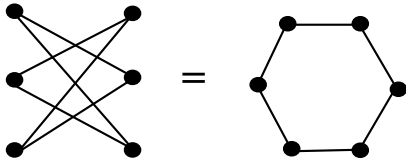
$$\sum_{i \in S} 2d_i(N - d_i) \geq \sum_{i \neq j \in S} d_i d_j \quad \text{for all subsets } S \subset \{1, \dots, K\}$$

$$K = 3$$

The threshold case for feasibility is

$$(d_1, d_2, d_3) = (d, d, N - d),$$

where  $d_1 \leq N/2$ .

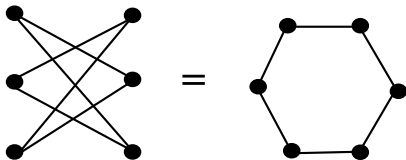


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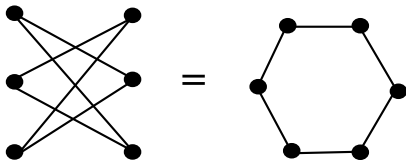
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- ▶ For dimension reasons, inclusions become **equalities**:

$$V_1 = U_3 = V_2 \subset U_1 = V_3 = U_2 \supset H_{21} V_1$$

## An eigenvector-like problem

Given generic  $N \times N$  matrix  $H$ , find

- ▶  $V \subset \mathbb{C}^N$ , subspace of dimension  $d$
- ▶  $U \subset \mathbb{C}^N$ , subspace of dimension  $e = N - d$

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- ▶ For  $e > d$ , variety of solutions of dimension

$$(N - (e - d))(e - d)$$



## Parametrizing the solution variety

Recall: Want to find  $V, U$  such that  $V \subset U$  and  $HV \subset U$ .

- ▶ Set  $\ell := \left\lfloor \frac{d}{e-d} \right\rfloor$
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$$U = S + T + \dots + H^{\ell-1}T$$

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Structure of whole variety seems complicated: when  $e = d + 1$ , then it is the **toric variety** for the Minkowski sum of hypersimplices  $\Delta_{e,N} + \Delta_{N,d}$ .

## Numbers of solutions

Return to

$$K \geq 3 \quad d = d_1 = \dots = d_k$$

Zero-dimensional when  $N = \frac{d(K+1)}{2}$ . The number of solutions is:

d	3	4	5	6	7
1	2	-	216	-	1,975,560
2	6	3700	388,407,960		
3	20	-		-	
4	70				
⋮	⋮				
d	$\binom{2d}{d}$				

## Number of solutions when $d = 1$

Assume  $d = d_1 = \cdots d_K = 1$  and  $2N = K + 1$ .

Degenerate each  $H_{ij}$  to a **rank 1** matrix:

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### Theorem

Number of solutions = number of **balanced orientations** of the graph  $G$

$G$  has edges  $t_j - s_i$  whenever  $i \neq j$ . Balanced orientation means that

$$\text{in degree}(v) = \text{out degree}(v) = \frac{K - 1}{2}$$

for all vertices  $v$  of  $G$ .

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- ▶ Can we **parametrize** the solution variety in more cases?
- ▶ What if the receivers and transmitters have **different numbers of antennas**?
- ▶ What if the channel matrices have the form

$$H_{ij} = \begin{bmatrix} \tilde{H}_{ij} & 0 & \cdots & 0 \\ 0 & \tilde{H}_{ij} & & 0 \\ \vdots & & \ddots & \vdots \\ 0 & 0 & \cdots & \tilde{H}_{ij} \end{bmatrix} ?$$